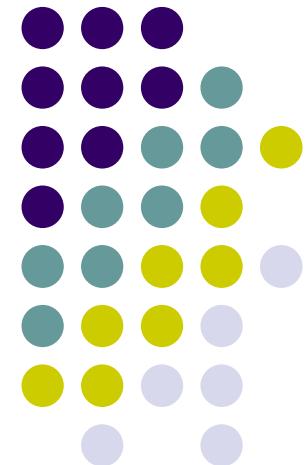


Midterm 1 Solutions

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I. Problem Formulation

I. A full adder inputs three bits (a, b, c_{in}) and outputs the sum and the carry bits (s, c_{out}). Write the **truth table** of the full adder and the corresponding **sum-of-products** canonical expressions of the two outputs.

Truth table:

a	b	c_{in}	s	c_{out}
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

SoP expressions:

$$\begin{aligned}s &= a'b'c_{in} + a'bc'_{in} + ab'c'_{in} + abc_{in} \\c_{out} &= abc'_{in} + ab'c_{in} + a'bc_{in} + abc_{in} \\&= ab + ac_{in} + bc_{in}\end{aligned}$$



II. Laws and Theorems of Boolean Algebra

II. Prove using Boolean algebra that $a'b' + a'c' + bc' = a'b' + bc'$.

Solution 1: use consensus

$$\begin{aligned} & a'b' + a'c' + bc' \\ &= a'b' + bc' + a'c' \quad \text{commutativity} \\ &= a'b' + bc' \quad \text{consensus} \end{aligned}$$

Solution 2: use combining & covering

$$\begin{aligned} & a'b' + a'c' + bc' \\ &= a'b' + (a'c'b' + a'c'b) + bc' \quad \text{combining} \\ &= (a'b' + a'b'c') + (bc'a' + bc') \quad \text{commutativity \& associativity} \\ &= a'b' + bc' \quad \text{covering} \end{aligned}$$



III. Laws and Theorems of Boolean Algebra

III. Prove using Boolean algebra that $(a' + b)(a' + c')(b' + c') = (a' + b)(b' + c')$.

Solution 1: use consensus

$$\begin{aligned} & (a' + b)(a' + c')(b' + c') \\ &= (a' + b)(b' + c')(a' + c') \quad \text{commutativity} \\ &= (a' + b)(b' + c') \quad \text{consensus} \end{aligned}$$

Solution 2: use combining & covering

$$\begin{aligned} & (a' + b)(a' + c')(b' + c') \\ &= (a' + b)(a' + c' + b)(a' + c' + b')(b' + c') \quad \text{combining} \\ &= ((a' + b)(a' + b + c'))((b' + c' + a')(b' + c')) \quad \text{commutativity \& associativity} \\ &= (a' + b)(b' + c') \quad \text{covering} \end{aligned}$$



IV. Karnaugh Map

IV. Use Karnaugh map to simplify function $f(a,b,c) = \sum m(1,4,7) + \sum d(3,6)$.

List all possible minimal two-level **sum-of-products** expressions. Show the Boolean expressions. No need for the logic diagram.

There are two forms of maps

F	a	b	c	00	01	11	10
	0	1	x	0	1	x	0
	1	1	0	1	x	1	x

PIs: $a'c, bc, ac', ab$

EPIs: $a'c, ac'$

F	c	a	b	00	01	11	10
	0	0	0	0	0	x	1
	1	1	0	x	1	1	0

All possible min SoP expressions:

$$f = a'c + ac' + bc$$

$$\text{and } f = a'c + ac' + ab$$



V. Karnaugh Map

V. Use Karnaugh map to simplify function $f(a,b,c) = \sum m(0,3,5) + \sum d(2,4)$.

List all possible minimal two-level **product-of-sums** expressions. Show the Boolean expressions. No need for the logic diagram.

There are two forms of maps

F	a	b	c	00	01	11	10
	1	bc		1	0	1	x
	x		a	x	1	0	0

F	c	a	b	00	01	11	10
	1	x	ab	1	x	0	x
	0	0		x	1	0	1

PIs: $a+b+c'$, $a'+b'$, $b'+c$, $a'+c$

EPIs: $a+b+c'$, $a'+b'$

All possible min PoS expressions:

$$f = (a+b+c')(a'+b')$$

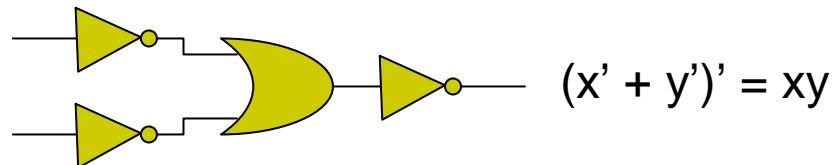
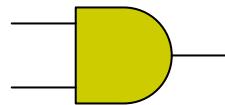


VI. Universal Gates

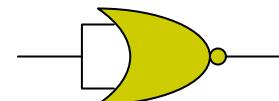
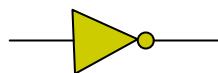
VI. Check if the set is universal and explain. Assume constants 0 and 1 are available as inputs.

i. {OR, NOT}

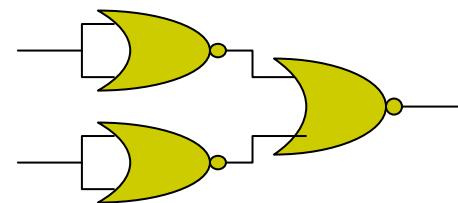
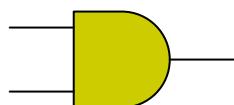
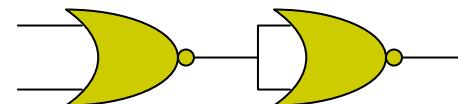
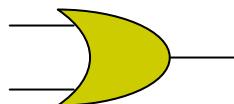
OR, NOT already exist, only need to construct AND



ii. {NOR}



$$(x + x)' = x'$$

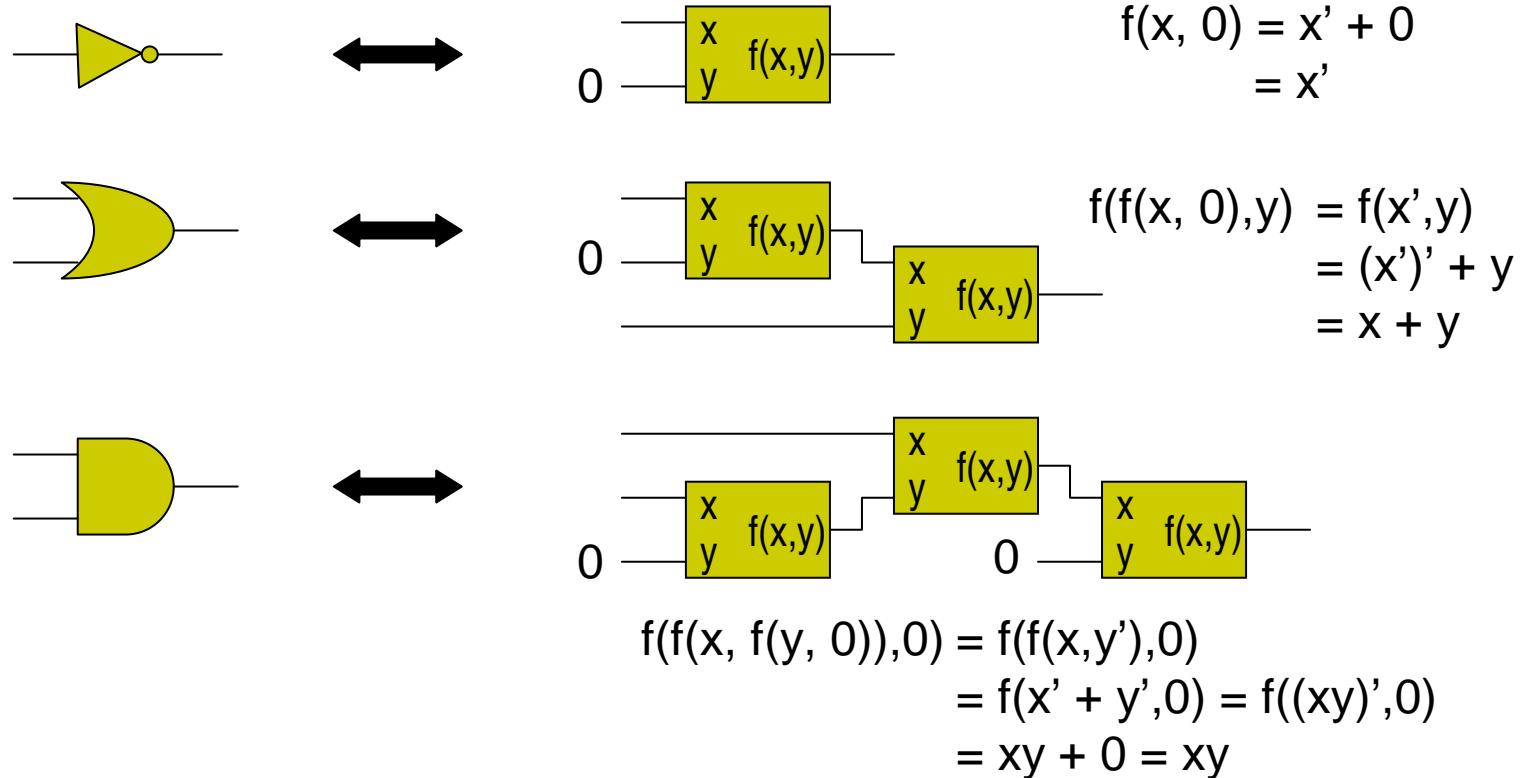




VI. Universal Gates

VI. Check if the set is universal and explain. Assume constants 0 and 1 are available as inputs.

iii. $\{f(x,y)\}$ where $f(x,y) = x' + y$

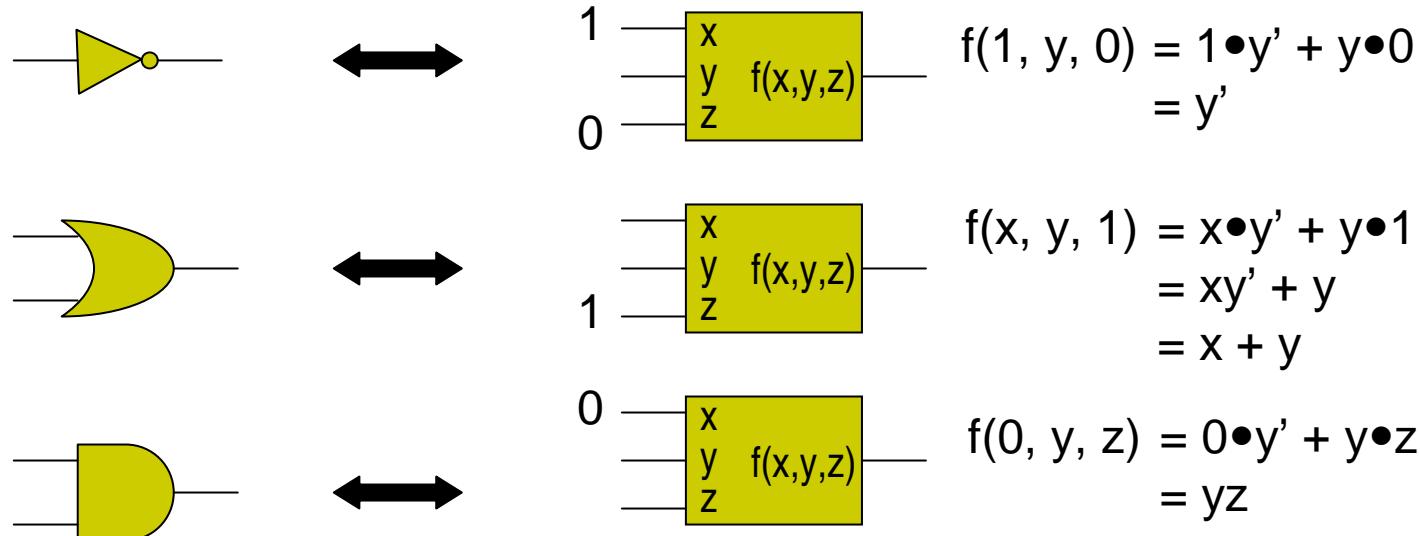




VI. Universal Gates

VI. Check if the set is universal and explain. Assume constants 0 and 1 are available as inputs.

iv. $\{f(x,y,z)\}$ where $f(x,y,z) = xy' + yz$





VII. Shannon's Expansion

VII. Describe the following function in minimal **sum-of-products** form

$$f(x, y) = xy' \oplus (x' + y) \oplus x'y \oplus y'$$

Expansion for variable x:

$$f(x, y) = x \bullet f(1, y) + x' \bullet f(0, y)$$

$$f(1, y) = y' \oplus y \oplus 0 \oplus y' = 1 \oplus 0 \oplus y' = 1 \oplus y' = y$$

$$f(0, y) = 0 \oplus 1 \oplus y \oplus y' = 1 \oplus y \oplus y' = y' \oplus y' = 0$$

$$f(x, y) = x \bullet y + x' \bullet 0 = xy$$

Expansion for variable y:

$$f(x, y) = y \bullet f(x, 1) + y' \bullet f(x, 0)$$

$$f(x, 1) = 0 \oplus 1 \oplus x' \oplus 0 = 1 \oplus x' \oplus 0 = x \oplus 0 = x$$

$$f(x, 0) = x \oplus x' \oplus 0 \oplus 1 = 1 \oplus 0 \oplus 1 = 1 \oplus 1 = 0$$

$$f(x, y) = x \bullet y + 0 \bullet y' = xy$$